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#### Attacking AUTOSAR using Software and Hardware Attacks

**Pascal Nasahl** Graz University of Technology

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  - Principal Security Analyst <a>D</a> Riscure

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Please visit Riscure's booth for more information! ③

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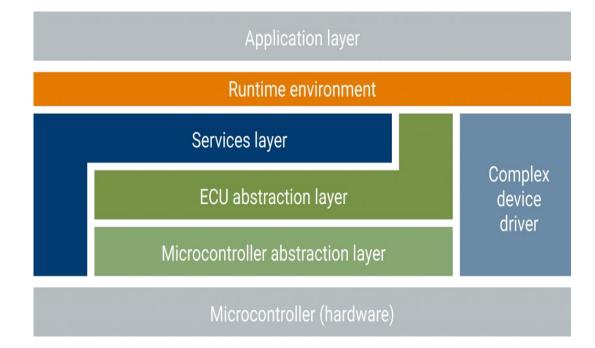
#### • Brief introduction to AUTOSAR Classic

- Brief introduction to AUTOSAR Classic
- Attacks on AUTOSAR

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- Case study

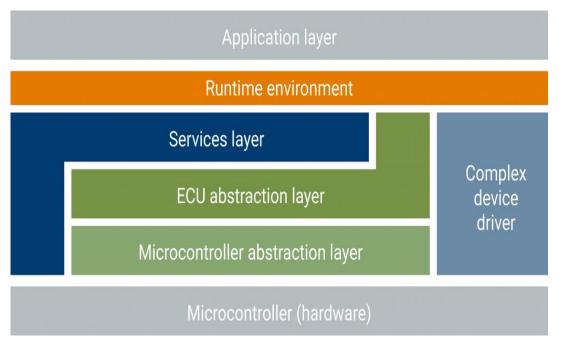
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- •Q&A



Layered software

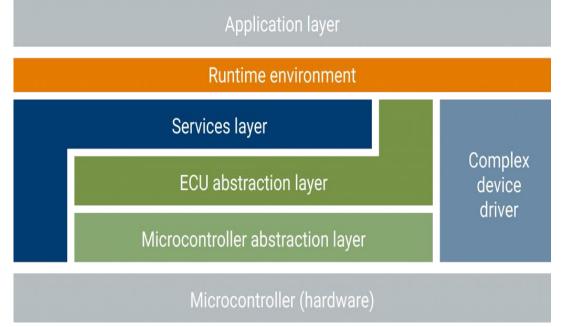
architecture



Layered software

architecture

 Most layers are independent from the Microcontroller

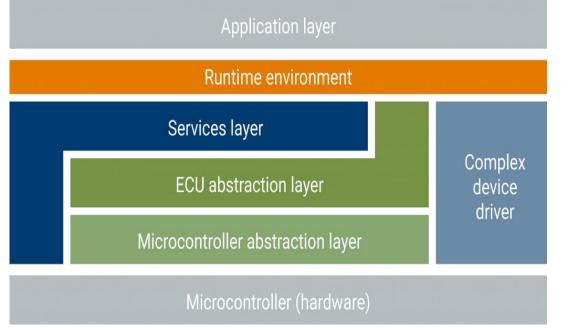


Layered software

architecture

 Most layers are independent from the Microcontroller

• Improve software reusability



| Application Layer |                                  |                                   |                                   |   |  |                             |                    |
|-------------------|----------------------------------|-----------------------------------|-----------------------------------|---|--|-----------------------------|--------------------|
|                   |                                  |                                   | Runtime E                         | nvironment                                  |  |                             |                    |
| Syste             | m Services                       | Memory<br>Services                | Crypto<br>Services                | Off-board<br>Communication<br>Services      | Communication<br>Services                | I/O Hardware<br>Abstraction | Complex<br>Drivers |
|                   | Onboard<br>Device<br>Abstraction | Memory<br>Hardware<br>Abstraction | Crypto<br>Hardware<br>Abstraction | Wireless<br>Communication<br>HW Abstraction | Communication<br>Hardware<br>Abstraction |                             |                    |
|                   | Microcontroller<br>Drivers       | Memory<br>Drivers                 | Crypto Drivers                    | Wireless<br>Communication<br>Drivers        | Communication<br>Drivers                 | I/O Drivers                 |                    |
| Microcontroller   |                                  |                                   |                                   |   |  |                             |                    |

| Application Layer |                                  |                                   |                                   |   |  |                             |                    |
|-------------------|----------------------------------|-----------------------------------|-----------------------------------|---|--|-----------------------------|--------------------|
|                   |                                  |                                   | Runtime E                         | nvironment                                  |  |                             |                    |
| Syste             | m Services                       | Memory<br>Services                | Crypto<br>Services                | Off-board<br>Communication<br>Services      | Communication<br>Services                | I/O Hardware<br>Abstraction | Complex<br>Drivers |
|                   | Onboard<br>Device<br>Abstraction | Memory<br>Hardware<br>Abstraction | Crypto<br>Hardware<br>Abstraction | Wireless<br>Communication<br>HW Abstraction | Communication<br>Hardware<br>Abstraction |                             |                    |
|                   | Microcontroller<br>Drivers       | Memory<br>Drivers                 | Crypto Drivers                    | Wireless<br>Communication<br>Drivers        | Communication<br>Drivers                 | I/O Drivers                 |                    |
| Microcontroller   |                                  |                                   |                                   |   |  |                             |                    |

#### Vulnerabilities can be introduced in any layer!

• Complex software; will contain bugs/vulnerabilities

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• Made by different vendors / developers

• Do you trust your suppliers?

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Mature code due to safety requirements
i.e. MISRA-C

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Mature code due to safety requirements
i.e. MISRA-C

#### Mature! But not guaranteed secure...

# What can go wrong?

### Potential MCAL vulnerabilities

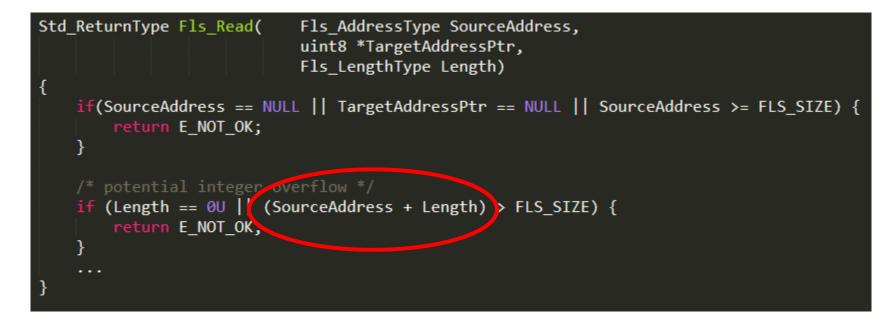
```
Std_ReturnType Fls_Read( Fls_AddressType SourceAddress,
uint8 *TargetAddressPtr,
Fls_LengthType Length)
```

```
if(SourceAddress == NULL || TargetAddressPtr == NULL || SourceAddress >= FLS_SIZE) {
    return E_NOT_OK;
}
```

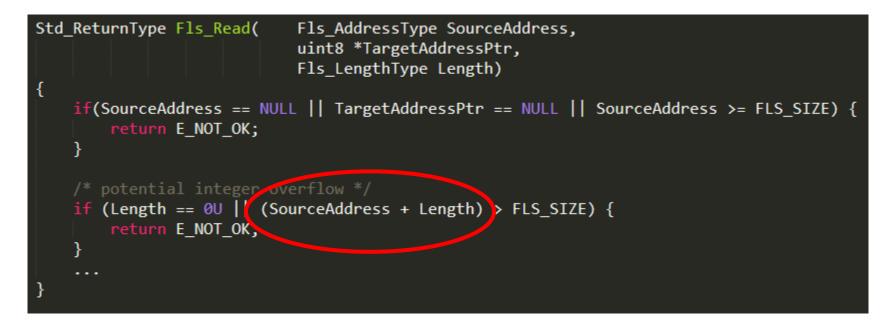
```
/* potential integer overflow */
if (Length == 0U || (SourceAddress + Length) > FLS_SIZE) {
    return E_NOT_OK;
}
...
```

{

### Potential MCAL vulnerabilities



## Potential MCAL vulnerabilities



#### Who verifies your MCAL for vulnerabilities?

#### What about MISRA-C?!

#### What about MISRA-C?!

| Dir 4.1                            | Run-time failures shall be minimized   | Dir 4.14   | The validity of values received from external sources shall be checked                           |
|------------------------------------|--|------------|--|
|                                    | C90 [Undefined 15, 19, 26, 30, 31, 32, 94]<br>C99 [Undefined 15, 16, 33, 40, 43–45, 48, 49, 113] |            | C90 [Undefined 15, 19, 26, 30, 31, 32, 94]<br>C99 [Undefined 15, 16, 33, 40, 43-45, 48, 49, 113] |
| CategoryRequiredApplies toC90, C99 |  | Category   | Required   |
|                                    | (30, (33   | Applies to | C90, C99   |

### What about MISRA-C?!

| Dir 4.1 | Run-time failures shall be minimized   | Dir 4.14   | The validity of values received from external sources shall be checked                           |
|---------|--|------------|--|
|         | C90 [Undefined 15, 19, 26, 30, 31, 32, 94]<br>C99 [Undefined 15, 16, 33, 40, 43–45, 48, 49, 113] |            | C90 [Undefined 15, 19, 26, 30, 31, 32, 94]<br>C99 [Undefined 15, 16, 33, 40, 43-45, 48, 49, 113] |
| 0,      | Required   | Category   | Required   |
|         | C90, C99   | Applies to | C90, C99   |

A directive is a guideline for which it is not possible to provide the full description necessary to perform a check for compliance. Additional information, such as might be provided in design documents or requirements specifications, is required in order to be able to perform the check. Static analysis tools may be able to assist in checking compliance with directives but different tools may place widely different interpretations on what constitutes a non-compliance.

#### You cannot conform to directives automagically...

#### What else?

### Vulnerabilities in complex software

## Vulnerabilities in complex software

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Ori Karliner @oriHCX · 18 Oct 2018

FreeRTOS TCP/IP Stack Vulnerabilities Put A Wide Range of Devices at Risk of Compromise: From Smart Homes to Critical Infrastructure Systems | Zimperium Mobile Security Blog



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## Who verifies your communication stack?

## Mitigating software vulnerabilities

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- Minimize the low hanging fruit
  - Secure coding standard, code checkers, ...

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  - Continuous security code reviews, ...
- Make it harder to exploit software vulnerabilities

• Software exploitation mitigations, ...

# Sufficiently complex software has vulnerabilities.

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Finding them is not always trivial...

Sufficiently complex software has vulnerabilities.

Finding them is not always trivial...

What if an attacker cannot find any or they are too difficult to exploit?

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• Attacker needs physcial access the ECU

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- Attacker often needs to open the ECU

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- Different types of HW attacks:
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- Attacker often needs to open the ECU
- Different types of HW attacks:
  - E.g. PCB-level, Fault injection, Side Channels, etc.
- Often a <u>stepping stone</u> for <u>more scalable attacks</u>

"Using FI to take control of an AUTOSAR-based ECU."

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- Demonstration ECU implemented using:
  - STM32F4 development board
  - Arctic Core for AUTOSAR v3.1

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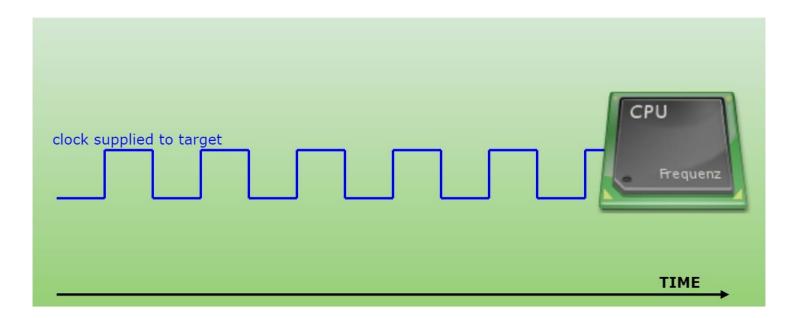
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- Attacking using a previously described FI fault model

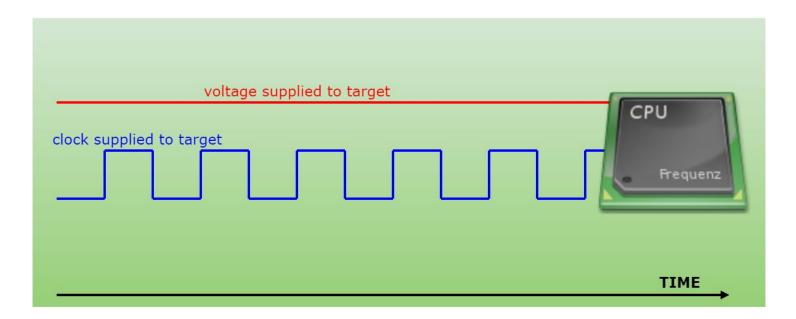
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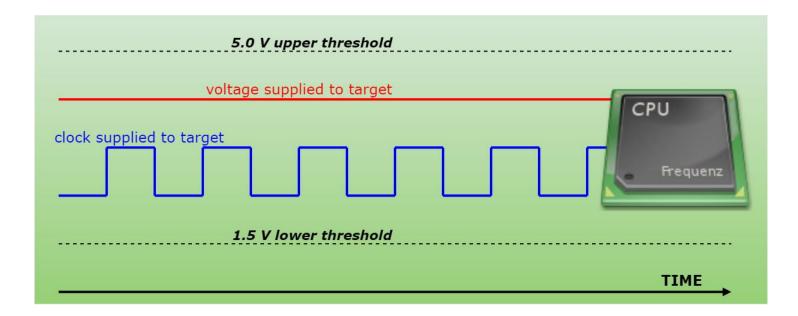
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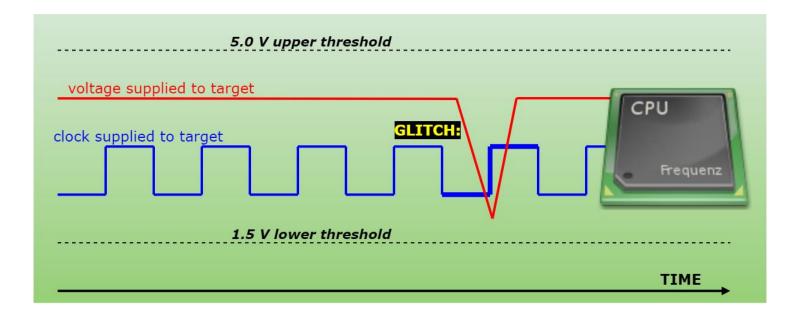
#### Fault Injection? Fault model?











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USB



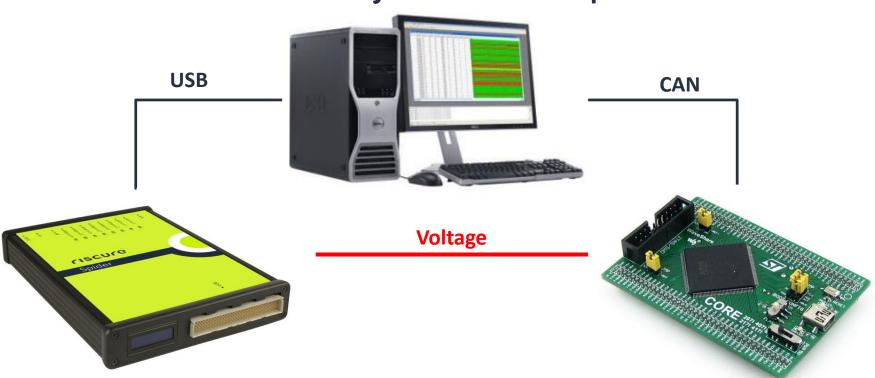


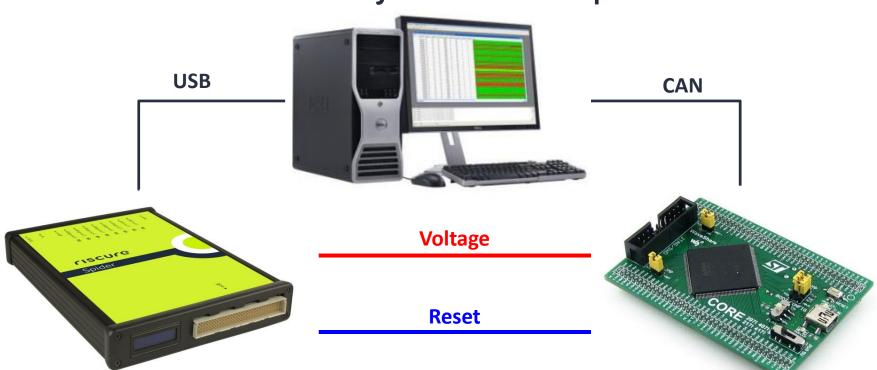
USB



CAN

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#### What can we do with fault injection?

"Instruction corruption."

"Instruction corruption."

• Glitches can be used to modify instruction

| Original instruction:              | Glitched instruction:              |
|------------------------------------|------------------------------------|
| add r0, r1, r3 1110 1011 0000 0001 | add r0, r1, r2 1110 1011 0000 0001 |
| 0000 0000 0000 0011                | 0000 0000 0000 0010                |

"Instruction corruption."

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|------------------------------------|------------------------------------|
| add r0, r1, r3 1110 1011 0000 0001 | add r0, r1, r2 1110 1011 0000 0001 |
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• In other words, we can modify software

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• In other words, we can modify software

• Fault injection breaks any software security model

### How can we use this to attack AUTOSAR-based ECUs?

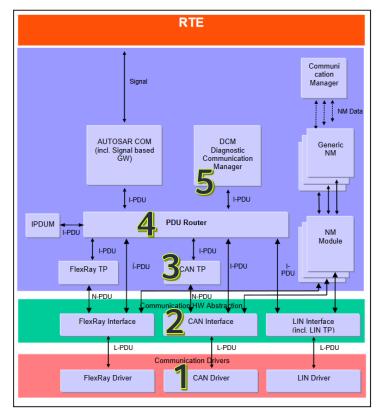
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#### • Our goal is to execute arbitrary code

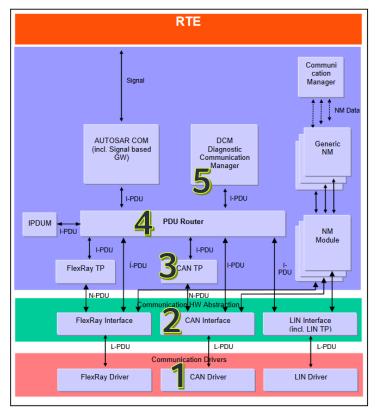
#### • Our goal is to execute arbitrary code

• Our only entry into the device is the CAN bus

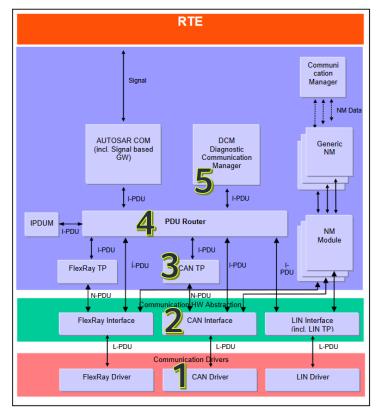
- Our goal is to execute arbitrary code
- Our only entry into the device is the CAN bus
- Of course, we have physical access...



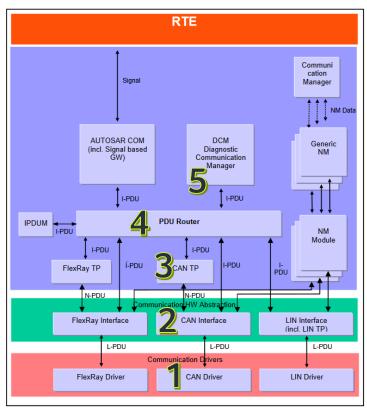
#### 1. CAN driver receives 8-byte CAN frame



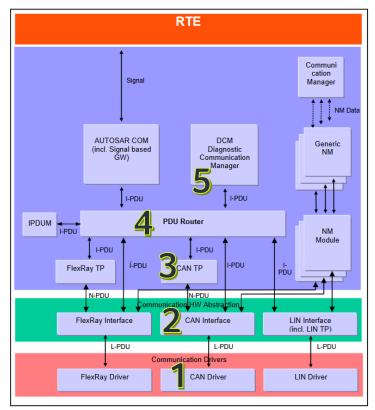
- 1. CAN driver receives 8-byte CAN frame
- 2. Frame passes the CAN interface



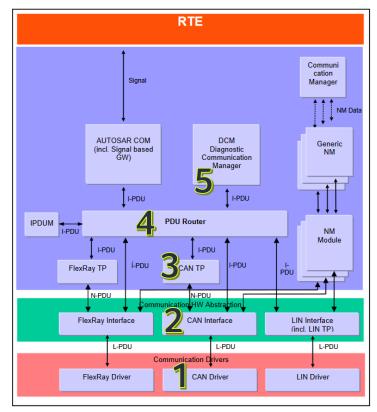
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- 1. CAN driver receives 8-byte CAN frame
- 2. Frame passes the CAN interface
- 3. Payload is reassembled by ISO-TP
- 4. Payload is copied to COM or DCM

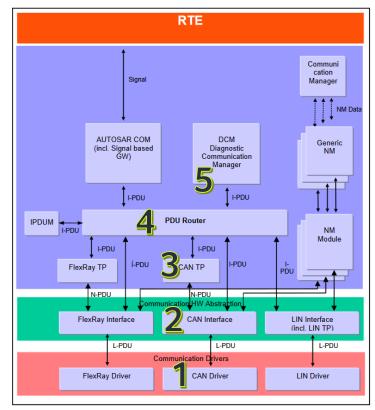


- 1. CAN driver receives 8-byte CAN frame
- 2. Frame passes the CAN interface
- 3. Payload is reassembled by ISO-TP
- 4. Payload is copied to COM or DCM
- 5. COM or DCM handles the payload



#### Where do we attack?!

- 1. CAN driver receives 8-byte CAN frame
- 2. Frame passes the CAN interface
- 3. Payload is reassembled by ISO-TP
- 4. Payload is copied to COM or DCM
- 5. COM or DCM handles the payload



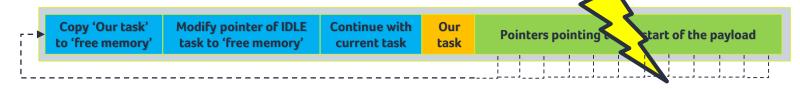
| Copy 'Our task'<br>to 'free memory' | Our<br>task |  |
|-------------------------------------|-------------|--|
|-------------------------------------|-------------|--|

|--|

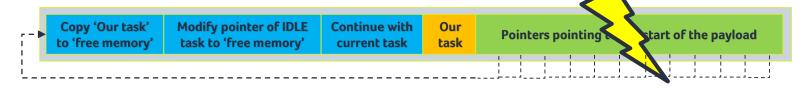
| Our<br>ask |  | Modify pointer of IDLE<br>task to 'free memory' |  |
|------------|--|---|--|
|------------|--|---|--|

| Modify pointer of IDLE<br>task to 'free memory' | Continue with<br>current task | Our<br>task | P   | oint | ers p | ointi | ing t | o th | e stai | rt of | the | payl | oad |
|---|-------------------------------|-------------|-----|------|-------|-------|-------|------|--------|-------|-----|------|-----|
|   |                               |             | 1   | 1    | 1     | 1     | 1     | 1    | 1      | 1     | 1   | 1    | 1.0 |
|   |                               |             |     | 1    | 1     | 1     | 1     | - 1  | 1      |       |     | 1    | 1   |
|   |                               |             | i i | 1    | 1     | 1     | 1     | 1    | 1      | 1     | 1   | 1    | 1   |
|   |                               |             |     |      |       |       |       |      |        |       |     |      |     |

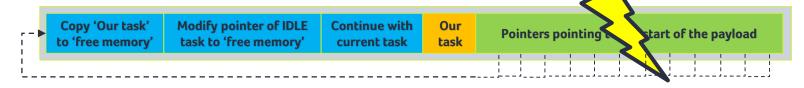
• Step 1: Send an ISO-TP CAN message (< 4096 bytes)



• Step 2: We inject the glitch when the pointers are being copied

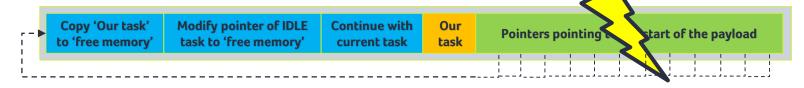


- Step 2: We inject the glitch when the pointers are being copied
- Step 3: Successful glitches load a pointer into the PC register



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- **Step 4:** MCU will execute the ISO-TP message (blue blocks)

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- Step 2: We inject the glitch when the pointers are being copied
- Step 3: Successful glitches load a pointer into the PC register
- **Step 4:** MCU will execute the ISO-TP message (blue blocks)
- Step 5: Wait for IDLE task to be scheduled and execute our task

## Why does this work?

BufReq\_ReturnType PduR\_CanTpCopyRxData(..., PduInfoType\* p, ...) {
 ...
 memcpy(p->SduDataPtr, source, p->SduLength);
 ...
}

```
BufReq_ReturnType PduR_CanTpCopyRxData(..., PduInfoType* p, ...) {
    ...
    memcpy(p->SduDataPtr, source, p->SduLength);
    ...
}
```

|                         | <memcpy>:</memcpy> |                                     | 1 | <br>r | г |  | r. |  |  |
|-------------------------|--------------------|-------------------------------------|---|-------|---|--|----|--|--|
|                         | mov                | ip, r0                              |   |       |   |  |    |  |  |
|                         | orr.w              | r3, r1, r0                          |   |       |   |  |    |  |  |
|                         | ands.w             | r3, r3, #3                          |   |       |   |  |    |  |  |
| <b>D</b> : <b>1 1 1</b> | bne.n              | 8008300 <memcpy+0xe8></memcpy+0xe8> |   |       |   |  |    |  |  |
| Disassembled            | subs               | r2, #64 ; 0x40                      |   |       |   |  |    |  |  |
|                         | n.33d              | 80082ac <memcpy+0x94></memcpy+0x94> |   |       |   |  |    |  |  |
| mamany()                | ldr.w              | <b>r3</b> , [r1], #4                |   |       |   |  |    |  |  |
| memcpy()                | str.w              | r3, [r0], #4                        |   |       |   |  |    |  |  |
|                         | subs               | r2, #64 ; 0x40                      |   |       |   |  |    |  |  |
|                         | bcs.n              | 8008228 <memcpy+0x10></memcpy+0x10> |   |       |   |  |    |  |  |
|                         | adds               | r2, #48 ; 0x30                      |   |       |   |  |    |  |  |
|                         | bcc.n              | 80082d4 <memcpy+0xbc></memcpy+0xbc> |   |       |   |  |    |  |  |
|                         |                    |                                     |   |       |   |  |    |  |  |

```
BufReq_ReturnType PduR_CanTpCopyRxData(..., PduInfoType* p, ...) {
    ...
    memcpy(p->SduDataPtr, source, p->SduLength);
    ...
}
```

|              | orr.w<br>ands.w             |
|--------------|-----------------------------|
| Disassembled | bne.n<br>subs               |
| Disassembled | bcc.n                       |
| memcpy()     | <mark>ldr.w</mark><br>str.w |
|              | subs                        |

| <memcpy>:</memcpy>                      |  |          | <memcpy>:</memcpy>                                   |  |
|---|--|----------|--|--|
| mov                                     | ip, r0   |          | mov  | ip, r0   |
| orr.w                                   | r3, r1, r0   |          | orr.w  | r3, r1, r0   |
| ands.w                                  | r3, r3, #3   |          | ands.w   | r3, r3, #3   |
| bne.n                                   | 8008300 <memcpy+0xe8></memcpy+0xe8>  |          | bne.n  | 8008300 <memcpy+0xe8></memcpy+0xe8>  |
| subs                                    | r2, #64 ; 0x40   |          | subs   | r2, #64 ; 0x40   |
| bcc.n                                   | 80082ac <memcpy+0x94></memcpy+0x94>  |          | bcc.n  | 80082ac <memcpy+0x94></memcpy+0x94>  |
| ldr.w                                   | <b>r3,</b> [r1], #4  |          | ldr.w  | <b>pc,</b> [r1], #4  |
| str.w                                   | r3, [r0], #4   |          | str.w  | r3, [r0], #4   |
| subs                                    | r2, #64 ; 0x40   |          | subs   | r2, #64 ; 0x40   |
| bcs.n                                   | 8008228 <memcpy+0x10></memcpy+0x10>  | <b>_</b> | bcs.n  | 8008228 <memcpy+0x10></memcpy+0x10>  |
| adds                                    | r2, #48 ; 0x30   |          | adds   | r2, #48 ; 0x30   |
| bcc.n                                   | 80082d4 <memcpy+0xbc></memcpy+0xbc>  |          | bcc.n  | 80082d4 <memcpy+0xbc></memcpy+0xbc>  |
| ldr.w<br>str.w<br>subs<br>bcs.n<br>adds | <b>r3</b> , [r1], #4<br>r3, [r0], #4<br>r2, #64 ; 0x40<br>8008228 <memcpy+0x10><br/>r2, #48 ; 0x30</memcpy+0x10> | 7        | <mark>ldr.w</mark><br>str.w<br>subs<br>bcs.n<br>adds | <pre>pc, [r1], #4 r3, [r0], #4 r2, #64 ; 0x40 8008228 <memcpy+0x1 #48="" 0x30<="" ;="" pre="" r2,=""></memcpy+0x1></pre> |

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BufReq_ReturnType PduR_CanTpCopyRxData(..., PduInfoType* p, ...) {
    ...
    memcpy(p->SduDataPtr, source, p->SduLength);
    ...
}
```

|              | <memcpy>:</memcpy> |                                     |   | <memcpy>:</memcpy> |                                     |
|--------------|--------------------|-------------------------------------|---|--------------------|-------------------------------------|
|              | mov                | ip, r0                              |   | mo∨                | ip, r0                              |
|              | orr.w              | r3, r1, r0                          |   | orr.w              | r3, r1, r0                          |
|              | ands.w             | r3, r3, #3                          |   | ands.w             | r3, r3, #3                          |
|              | bne.n              | 8008300 <memcpy+0xe8></memcpy+0xe8> |   | bne.n              | 8008300 <memcpy+0xe8></memcpy+0xe8> |
| Disassembled | subs               | r2, #64 ; 0x40                      |   | subs               | r2, #64 ; 0x40                      |
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| momcnv()     | ldr.w              | <b>r3</b> , [r1], #4                |   | <mark>ldr.w</mark> | <b>pc,</b> [r1], #4                 |
| memcpy()     | str.w              | r3, [r0], #4                        |   | str.w              | r3, [r0], #4                        |
|              | subs               | r2, #64 ; 0x40                      |   | subs               | r2, #64 ; 0x40                      |
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|              | adds               | r2, #48 ; 0x30                      |   | adds               | r2, #48 ; 0x30                      |
|              | bcc.n              | 80082d4 <memcpy+0xbc></memcpy+0xbc> |   | bcc.n              | 80082d4 <memcpy+0xbc></memcpy+0xbc> |

We take control of the Program Counter (PC) during the copy!

#### We have our own task. Now what?!

• Extract information (secrets)

- Extract information (secrets)
- Analyze firmware dynamically

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- Analyze firmware dynamically
- Perform additional attacks (e.g. side channel attack)

- Extract information (secrets)
- Analyze firmware dynamically
- Perform additional attacks (e.g. side channel attack)
- Add (malicious) and/or change functionality

#### Is all hope lost?



• Adhere to (automotive) security guidelines/standards

- Adhere to (automotive) security guidelines/standards
- Make use of strong (hardware-based) security

- Adhere to (automotive) security guidelines/standards
- Make use of strong (hardware-based) security
- Minimize attack surface and increase attack complexity

- Adhere to (automotive) security guidelines/standards
- Make use of strong (hardware-based) security
- Minimize attack surface and increase attack complexity
- Consult internal/external embedded security experts

To wrap up...





#### • Devices (incl. AUTOSAR-based ECUs) will be hacked

#### Takeaways

# Devices (incl. AUTOSAR-based ECUs) will be hacked Not AUTOSAR's fault!

#### Takeaways

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- No (known) software vulnerabilities ≠ secure
- Hardware attacks are efficient and do scale

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#### Thank you. Questions?

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